GIFT-COFB Final Round Updates

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1 Third-party security analysis

Regarding GIFT-128 block cipher, third-party security evaluations have continued to be published since the last revision. With respect to linear attacks, the number of attacked rounds has increased significantly, but this does not change the fact that differential cryptanalysis remains more effective against GIFT-128. These attacks were discovered using automated evaluation tools and research on more efficient tools for the GIFT network have also been presented [20]. No progress has been made on differential cryptanalysis, and no new (unknown) cryptanalysis method has been discovered. This is a strong evidence of the reliability of GIFT-128. We give below some comments on a few references.

- Sun et al. [26] and its addendum [27] detected a 25-round linear cryptanalysis against GIFT-128. An attack taking into account the AEAD restrictions was also proposed, which recovers the key of GIFT-COFB if the primitive is reduced to 16 rounds. Note that differential cryptanalysis still works better on GIFT for the moment.
- Cui *et al.* proposed a 20-round linear cryptanalysis against GIFT-128 [10], which only works for a smaller number of rounds than previous results.
- The work by Khalesi and Ahmadian searches for minimum data complexity of the integral distinguisher [19]. Regarding GIFT-128, it confirms that the previously known result is actually the best.
- The paper by Hu *et al.* [12] proves the non-existence of impossible differential with one active superbox in both ends for 8-round reduced GIFT-128.
- The paper by Anubhab Baksi [3] presents the optimal linear bounds for 11 and 12 rounds of GIFT-128, extending from the best-known result on 10rounds.

We recall that GIFT-128 has 40 rounds, while the best known attack can only reach 27 rounds (which does not apply to GIFT-COFB due to the data limitation, etc.). This leaves a very ample security margin.

The provable security aspects of GIFT-COFB and its variants have been studied and updated since the last revision of our specification document. We list the relevant works below.

- Khairallah [18] (ToSC 2022) presented an attack against a version of COFB mode shown in Journal of Cryptology [8]. Due to the difference in the modes, the presented attack is not applicable to GIFT-COFB.
- As an internal security evaluation of GIFT-COFB, Inoue, Iwata and Minematsu (IIM22, ACNS 2022 and NIST LWC workshop 2022) [14, 15] showed an attack that has a higher success probability than claimed. Their attack does not break the claimed bit security. A revised version of the proof (and the bound) maintaining the original bit security was presented in [5] (see the next section).
- Liu et al. [22] presented fault/side-channel attacks against unprotected implementation of GIFT-COFB.
- Reshma *et al.* [25] showed neural network-based distinguishers on reduced GIFT-COFB. The attack is very weak, as it only distinguishes between 2 and 6 rounds of GIFT-COFB's ciphertext from random data.

2 New security proofs/arguments

GIFT-COFB team revised the security proof of the mode and updated the bound by incorporating the fix suggested by IIM22. The fixed bound maintains the original bit security. The updated specification document (version 1.2) is posted at NIST LWC mailing list in May 2022⁸. The revised proof and the bound are shown in [5].

As an additional security feature, Inoue *et al.* [13] showed that GIFT-COFB has <u>resilience</u> against nonce misuse/repeat, which is a relaxed notion from the nonce misuse resistance introduced by Ashur *et al.* [2]. The presented bound is n/4 bits for n = 128, so not strong enough in general, but at least it does not lose everything in the event of nonce misuse. In contrast, the current NIST standard GCM has no resilience against nonce misuse. This work also supports the correctness of the design.

In summary, we actively reviewed the provable security of GIFT-COFB to increase the confidence in its security. As a result, we maintain the originally claimed bit security in a more transparent way.

⁸ https://groups.google.com/a/list.nist.gov/g/lwc-forum/c/7BmjTeE-NsY

Table 1. Summary of third-party analysis result on GIFT. Rounds with asterisk (*)
are optimal results. SK – single-key, RK – related-key, LC – linear cryptanalysis, DC
– differential cryptanalysis.

Distinguisher SK 11 Integral 1 - 2 ¹²⁷ - [11] SK 11* Integral 1 - 2 ¹²⁷ - [19] SK 9* LC 2 ⁻⁴⁴ - - - [16] SK 10* LC 2 ⁻⁵² - - [29] SK 15 LC 2 ^{-117,43} - - [29] SK 19 LC 2 ^{-117,43} - - [26] SK 19 LC 2 ^{-117,43} - - [23] SK 10* DC 2 ^{-49,4} - - [23] SK 11* DC 2 ^{-51,4} - - [23] SK 13* DC 2 ^{-67,8} - - [23] SK 13* DC 2 ^{-90,415} - - [16] SK 16* DC 2 ^{-90,415}	Setting	Rounds	Approach	Prob.	Time	Data	Mem.	Ref.						
SK 11 Integral 1 - 2^{127} - [11] SK 9* LC 2^{-44} - - - [16] SK 10* LC 2^{-52} - - - [16] SK 15 LC 2^{-109} - - - [29] SK 16 LC $2^{-117.43}$ - - [27] SK 19 LC $2^{-117.43}$ - - [27] SK 19 LC $2^{-123.11}$ - - [27] SK 19* DC $2^{-49.4}$ - - [23] SK 11* DC $2^{-67.8}$ - - [23] SK 13* DC $2^{-67.8}$ - - [16] SK 14* DC $2^{-90.415}$ - - [16] SK 16* DC $2^{-10.435}$	Distinguisher													
SK 11* Integral 1 - 2^{127} - [19] SK 9* LC 2^{-44} - - - [16] SK 10* LC 2^{-52} - - - [16] SK 16 LC 2^{-107} - - [27] SK 19 LC $2^{-117.43}$ - - [23] SK 19* DC $2^{-45.4}$ - - [23] SK 10* DC $2^{-49.4}$ - - [23] SK 11* DC $2^{-67.8}$ - - [23] SK 13* DC $2^{-67.8}$ - - [16] SK 16* DC $2^{-90.415}$ - - [16] SK 16* DC $2^{-10.435}$ - - [16] SK 18* DC $2^{-103.415}$ -	SK	11	Integral	1	-	2^{127}	-	[11]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	SK	11*	Integral	1	-	2^{127}	-	[19]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	SK	9*	LC	2^{-44}	-	-	-	[16]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	SK	10*	LC	2^{-52}	-	-	-	[16]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	SK	15	LC	2^{-109}	-	-	-	[29]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	SK	16	LC	2^{-122}	-	-	-	[10]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	SK	19	LC	$2^{-117.43}$	-	-	-	[26]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	SK	19	LC	$2^{-123.11}$	-	-	-	[27]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	SK	9*	DC	$2^{-45.4}$	-	-	-	[23]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	SK	10*	DC	$2^{-49.4}$	-	-	-	[23]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	SK	11*	DC	$2^{-54.4}$	-	-	-	[23]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	SK	12*	DC	$2^{-60.4}$	-	-	-	[23]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	SK	13*	DC	$2^{-67.8}$	-	-	-	[23]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	SK	14*	DC	$2^{-79.000}$	-	-	-	[16]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	SK	15*	DC	$2^{-85.415}$	-	-	-	[16]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	SK	16*	DC	$2^{-90.415}$	-	-	-	[16]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	SK	17*	DC	$2^{-96.415}$	-	-	-	[16]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	SK	18	DC	2^{-109}	-	-	-	[28]						
$\begin{tabular}{ c c c c c c c c c c c c c c c c c c c$	SK	18*	DC	$2^{-103.415}$	-	-	-	[16]						
$\begin{tabular}{ c c c c c c c c c c c c c c c c c c c$	SK	19	DC	$2^{-110.83}$	-	-	-	[16]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	SK	20	DC	$2^{-121.415}$	-	-	-	[21]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	SK	20	DC	$2^{-120.245}$	-	-	-	[17]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	SK	20	DC	$2^{-121.813}$	-	-	-	[29]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	SK	21	DC	$2^{-126.4}$	-	-	-	[23]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	RK	7	DC	$2^{-15.83}$	-	-	-	[7]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	RK	10	DC	$2^{-72.66}$	-	-	-	[7]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	RK	19	Boomerang	$2^{-121.2}$	-	-	-	[24]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	RK	19	Boomerang	$2^{-109.626}$	-	-	-	[17]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$			Ke	y-Recover	y.									
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	SK	20	LC	-	$2^{112.28}$	2^{126}	2^{65}	[10]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	SK	22	LC	-	2^{117}	2^{117}	2^{78}	[29]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	SK	24	LC	-	$2^{124.45}$	$2^{122.55}$	2^{105}	[26]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	SK	25	LC	-	$2^{124.75}$	$2^{126.77}$	2^{96}	[27]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	SK	22	DC	-	2^{114}	2^{114}	2^{53}	[28]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	SK	26	DC	-	$2^{124.415}$	2^{109}	2^{109}	[21]						
$\begin{tabular}{ c c c c c c c c c c c c c c c c c c c$	SK	26	DC	-	$2^{123.245}$	$2^{123.245}$	2^{109}	[17]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	SK	27	DC	-	$2^{124.83}$	$2^{123.53}$	2^{80}	[29]						
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$	RK	21	Boomerang	-	$2^{126.6}$	$2^{126.6}$	$2^{126.6}$	[24]						
RK 23 Rectangle - $2^{126.89}$ $2^{121.31}$ $2^{121.63}$ [17]	RK	22	Boomerang	-	$2^{112.63}$	$2^{112.63}$	2^{52}	[17]						
	RK	23	Rectangle	-	$2^{126.89}$	$2^{121.31}$	$2^{121.63}$	[17]						

ach Prob. Time Data Mem Ref Setting Box inds Appr

3 New software and hardware implementations

3.1 Bit serial implementation of GIFT-COFB

It is a well-known fact that implementing a hardware block cipher in a bitserial manner, which advances only one bit in the computation pipeline in each clock cycle, results in the smallest circuits. Nevertheless, an efficient bit-serial circuit for a mode of operation that utilizes finite field arithmetic with multiple constants has yet to be demonstrated in the literature.

In a recent work [6], this issue regarding efficient field arithmetic in bit- serial circuits has been addressed. As a result lightweight circuit for GIFT-COFB is proposed, that occupies less than 1500 GE, making it the to-date most area-efficient implementation of this construction. In a second step, it is demonstrated how the additional operations in the mode can be executed concurrently with block cipher operations itself so that the total latency is significantly reduced whilst incurring only a modest area increase. Finally, a first-order threshold implementation of GIFT-COFB is proposed, in which first-order side-channel security is experimentally verified.

Summary: A total of 3 bit-serial circuits are proposed that stand as the to-date most area-efficient GIFT-COFB implementations known in the literature.

- 1. GIFT-COFB-SER-S: This circuit represents an effective transformation of the <u>swap-and-rotate</u> GIFT-128 scheme into the GIFT-COFB mode of operation minimizing its area footprint.
- 2. GIFT-COFB-SER-F: Subsequently, it was observed that the interspersing of block cipher invocations with calls to the finite field module as found in the baseline GIFT-COFB design can be reordered by leveraging its inherent mathematical structure in order to further optimize the overall latency of GIFT-COFB-SER-S while only incurring a modest area increase.
- 3. GIFT-COFB-SER-TI: A bit-serial first-order threshold implementation is proposed based on GIFT-COFB-SER-F whose security is experimentally verified through statistical tests on signal traces obtained by measuring the implemented circuit on a SAKURA-G side-channel evaluation FPGA board.
- 4. All of the proposed schemes are synthesized on ASIC platforms using multiple standard cell libraries and compared with results to existing bit-serial implementations of NIST LWC candidate submissions, indicating the designs are among the smallest currently in the competition. A brief overview of the synthesis results is tabulated in Table 2.

Moreover, we are also aware of the following implementations:

- a protected implementation of GIFT-COFB [1]
- a white-box implementations of GIFT-128 [9]

	Datapath	Area	Latency	Power	Energy	Reference
	Bits	GE	Cycles	μW	nJ	
SUNDAE-GIFT	1	1201	92544	55.48	513.4	[4]
SAEAES	1	1350	24448	84.47	206.5	[4]
Romulus	1	1778	55431	82.28	456.1	[4]
SKINNY-AEAD	1	3589	72960	143.7	1048	[4]
GIFT-COFB	128	3927	400	156.3	6.254	[5]
GIFT-COFB-SER-S	1	1443	54784	50.11	275.8	[6]
GIFT-COFB-SER-F	1	1485	51328	62.15	319.8	[6]
GIFT-COFB-SER-TI	1	3384	51328	158.1	813.5	[6]

Table 2. Synthesis results overview for lightweight block cipher based NIST LWC competitors using the STM 90 nm cell library at a clock frequency of 10 MHz. Latency and energy correspond to the encryption of 128 bits of AD and 1024 message bits. Highlighted schemes are NIST LWC finalists.

4 Target applications

GIFT-COFB performs extremely well in all hardware situations (low area, low energy, low power), and also very well in software due to its fixsliced representation.

5 Others

Akiko Inoue (NEC) joined the team of GIFT-COFB.

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